



A NEW LOOK

at the

STATE REPRESENTATION PROBLEM

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RUG

OUTLINE

- 1. Behavioral systems**
- 2. State systems**
- 3. Linear differential systems**
- 4. State representations**
- 5. State construction**
- 6. Balanced realizations**
- 7. State for PDE's**
- 8. Conclusions**

BEHAVIORAL SYSTEMS

THE BASIC CONCEPT

A dynamical system = $\Sigma = (\mathbb{T}, \mathbb{W}, \mathcal{B})$

$\mathbb{T} \subseteq \mathbb{R}$, the time-axis (= the relevant time instances),

\mathbb{W} , the signal space (= where the variables take on their values),

$\mathcal{B} \subseteq \mathbb{W}^{\mathbb{T}}$: the behavior (= the admissible trajectories).

$$\Sigma = (\mathbb{T}, \mathbb{W}, \mathfrak{B})$$

For a trajectory $w : \mathbb{T} \rightarrow \mathbb{W}$, we thus have:

$w \in \mathfrak{B}$: the model **allows** the trajectory w ,
 $w \notin \mathfrak{B}$: the model **forbids** the trajectory w .

Usually, $\mathbb{T} = \mathbb{R}$, or $[0, \infty)$ (in continuous-time systems),
or \mathbb{Z} , or \mathbb{N} (in discrete-time systems).

Usually, $\mathbb{W} \subseteq \mathbb{R}^w$ (in lumped systems);

a function space

(in distributed systems, with time a distinguished variable);

or a finite set (in DES).

Emphasis today: $\mathbb{T} = \mathbb{R}$, $\mathbb{W} = \mathbb{R}^w$,

$\mathfrak{B} =$ solutions of system of linear constant coefficient ODE's.

EXAMPLES

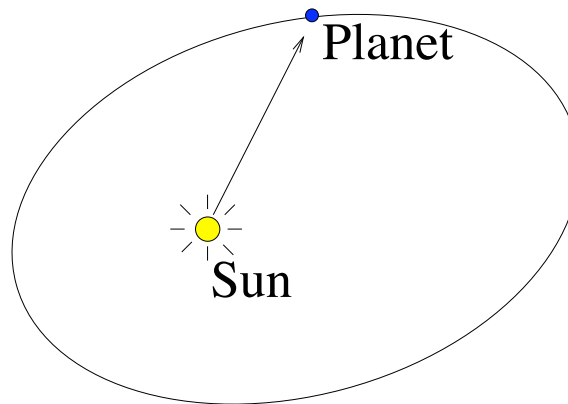
1. Planetary orbits

$T = \mathbb{R}$ (time),

$W = \mathbb{R}^3$ (position),

$\mathfrak{B} =$ planetary orbits \cong Kepler's laws:

ellipses, = areas in = time, $\frac{(\text{period})^2}{(\text{axis})^3} = \text{constant}$.



Planetary orbits

2. Input / output systems

$$\begin{aligned} f_1(y(t), \frac{d}{dt}y(t), \frac{d^2}{dt^2}y(t), \dots, t) \\ = f_2(u(t), \frac{d}{dt}u(t), \frac{d^2}{dt^2}u(t), \dots, t) \end{aligned}$$

$\mathbb{T} = \mathbb{R}$ (time),

$\mathbb{W} = \mathbb{U} \times \mathbb{Y}$ (input \times output signal spaces),

$\mathfrak{B} =$ **all input / output pairs.**

3. Codes

\mathbb{A} = the code alphabet, say, $\mathbb{A} = \mathbb{F}^w$, \mathbb{F} a finite field,

\mathbb{I} = an index set, say,

$\mathbb{I} = (1, \dots, n)$ in block codes,

$\mathbb{I} = \mathbb{N}$ or \mathbb{Z} in convolutional codes,

$\mathcal{C} \subseteq \mathbb{A}^{\mathbb{I}}$ = **the code**; yields the system $\Sigma = (\mathbb{I}, \mathbb{A}, \mathcal{C})$.

Redundancy structure, error correction possibilities, etc., are visible in the code behavior \mathcal{C} .

It is the central object of study.

Encoder & decoder can be put (temporarily) into the background.

4. Formal languages

\mathbb{A} = a (finite) alphabet,

$\mathcal{L} \subseteq \mathbb{A}^*$ = **the language** = all 'legal' 'words' $a_1 a_2 \cdots a_k \cdots$
yields the system $\Sigma = (\mathbb{N}, \mathbb{A}, \mathcal{L})$.

\mathbb{A}^* = all finite strings with symbols from \mathbb{A} .

Examples: All words appearing in the *van Dale*
All \LaTeX documents

LATENT VARIABLE SYSTEMS

First principles models \rightsquigarrow

A dynamical system with latent variables = $\Sigma_L = (\mathbb{T}, \mathbb{W}, \mathbb{L}, \mathcal{B}_{\text{full}})$

$\mathbb{T} \subseteq \mathbb{R}$, the *time-axis* (= the set of relevant time instances).

\mathbb{W} , the *signal space* (= the variables that the model aims at).

\mathbb{L} , the *latent variable space* (= the **auxiliary** modeling variables).

$\mathcal{B}_{\text{full}} \subseteq (\mathbb{W} \times \mathbb{L})^{\mathbb{T}}$: the full behavior

(= the pairs $(w, \ell) : \mathbb{T} \rightarrow \mathbb{W} \times \mathbb{L}$ that the model declares possible).

THE MANIFEST BEHAVIOR

Call the elements of \mathbb{W} *'manifest' variables*,
those of \mathbb{L} *'latent' variables*.

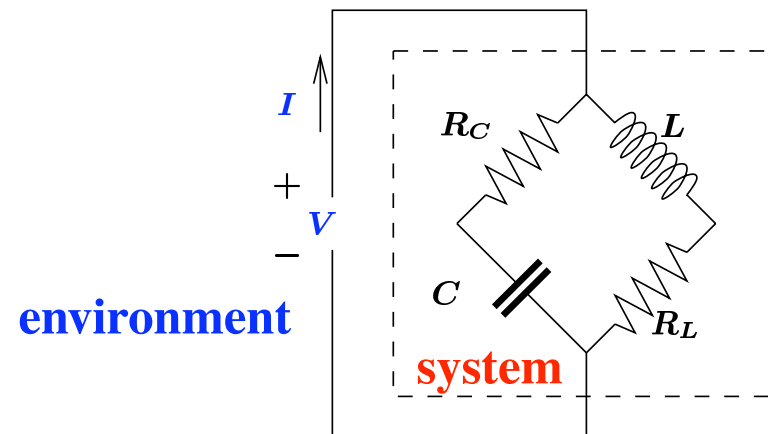
The latent variable system $\Sigma_L = (\mathbb{T}, \mathbb{W}, \mathbb{L}, \mathfrak{B}_{\text{full}})$ induces
the *manifest system* $\Sigma = (\mathbb{T}, \mathbb{W}, \mathfrak{B})$ with *manifest behavior*

$$\mathfrak{B} = \{w : \mathbb{T} \rightarrow \mathbb{W} \mid \exists \ell : \mathbb{T} \rightarrow \mathbb{L} \text{ such that } (w, \ell) \in \mathfrak{B}_{\text{full}}\}$$

In convenient equations for \mathfrak{B} , the latent variables are *'eliminated'*.

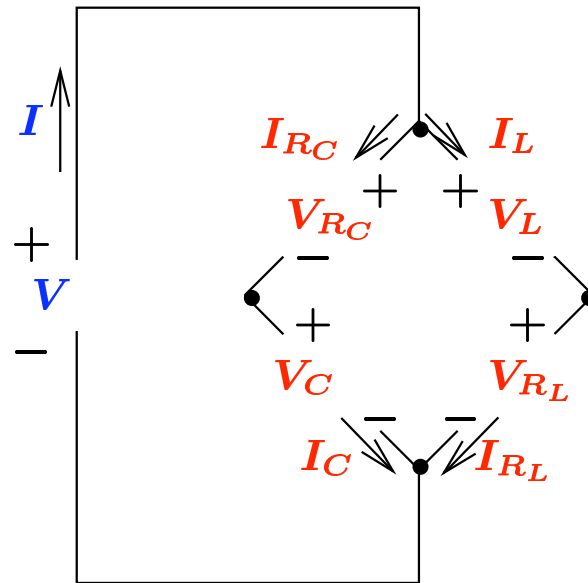
EXAMPLES

1. RLC - circuit



RLC - circuit

!! Model the relation between V and I !!



The circuit graph

SYSTEM EQUATIONS

Introduce the following additional variables:

the **voltage across** and the **current in** each branch:

$$V_{R_C}, I_{R_C}, V_C, I_C, V_{R_L}, I_{R_L}, V_L, I_L.$$

Constitutive equations (CE):

$$V_{R_C} = R_C I_{R_C}, \quad V_{R_L} = R_L I_{R_L}, \quad C \frac{d}{dt} V_C = I_C, \quad L \frac{d}{dt} I_L = V_L$$

Kirchhoff's voltage laws (KVL):

$$V = V_{R_C} + V_C, \quad V = V_L + V_{R_L}, \quad V_{R_C} + V_C = V_L + V_{R_L}$$

Kirchhoff's current laws (KCL):

$$I = I_{R_C} + I_L, \quad I_{R_C} = I_C, \quad I_L = I_{R_L}, \quad I_C + I_{R_L} = I$$

Formalization as a latent variable system:

$$\mathbb{T} = \mathbb{R},$$

$\mathbb{W} = \mathbb{R}^2$ - manifest variables: the **port voltage and current**,

$\mathbb{L} = \mathbb{R}^8$ - latent variables: the **branch voltages and currents**,

$\mathfrak{B}_{\text{full}} =$ all functions $(V, I, V_{RC}, I_{RC}, V_C, I_C, V_{RL}, I_{RL}, V_L, I_L)$
that satisfy the CE's, KCL, and KVL,

$\mathfrak{B} =$ the functions (V, I) that satisfy the 'eliminated' port
equations.

RELATION BETWEEN V and I

After some calculations ‘**elimination**’, we obtain the **port equations**:

Case 1: $CR_C \neq \frac{L}{R_L}$.

$$\left(\frac{R_C}{R_L} + \left(1 + \frac{R_C}{R_L}\right)CR_C \frac{d}{dt} + CR_C \frac{L}{R_L} \frac{d^2}{dt^2}\right)V = \left(1 + CR_C \frac{d}{dt}\right)\left(1 + \frac{L}{R_L} \frac{d}{dt}\right)R_C I.$$

Case 2: $CR_C = \frac{L}{R_L}$.

$$\left(\frac{R_C}{R_L} + CR_C \frac{d}{dt}\right)V = \left(1 + CR_C \frac{d}{dt}\right)R_C I$$

These are the **exact** relations between V and I !

2. Input /state / output systems

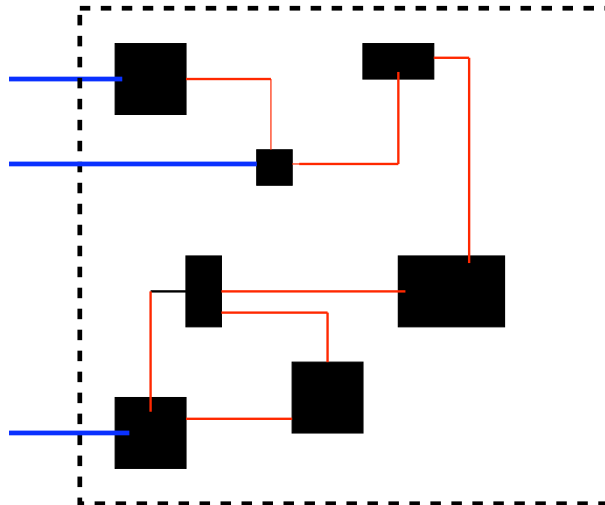
$$\frac{d}{dt}\mathbf{x}(t) = \mathbf{f}(\mathbf{x}(t), \mathbf{u}(t)); \quad \mathbf{y}(t) = \mathbf{h}(\mathbf{x}(t), \mathbf{u}(t)),$$

$$\mathbb{T} = \mathbb{R}, \mathbb{W} = \mathbb{U} \times \mathbb{Y}, \mathbb{L} = \mathbb{X},$$

$\mathfrak{B}_{\text{full}} = \text{all } (\mathbf{u}, \mathbf{y}, \mathbf{x}) : \mathbb{R} \rightarrow \mathbb{U} \times \mathbb{Y} \times \mathbb{X} \text{ that satisfy these equations,}$

$\mathfrak{B} = \text{all (input / output)-pairs.}$

3. Interconnected systems \cong 'First principles' models



Interconnected system

External variables \rightsquigarrow **manifest;**

Internal variables \rightsquigarrow **latent.**

4. Grammars

A convenient way to specify a formal language, whose essence is captured by **latent** variables, is through *grammars*.

STATE SYSTEMS

THE NOTION OF STATE

The **latent variable system**

$$\Sigma_X = (\mathbb{T}, \mathbb{W}, \mathbb{X}, \mathfrak{B}_{\text{full}})$$

is said to be a **state system** if

$$(w_1, x_1), (w_2, x_2) \in \mathfrak{B}_{\text{full}} \quad \text{and} \quad x_1(t_0) = x_2(t_0)$$

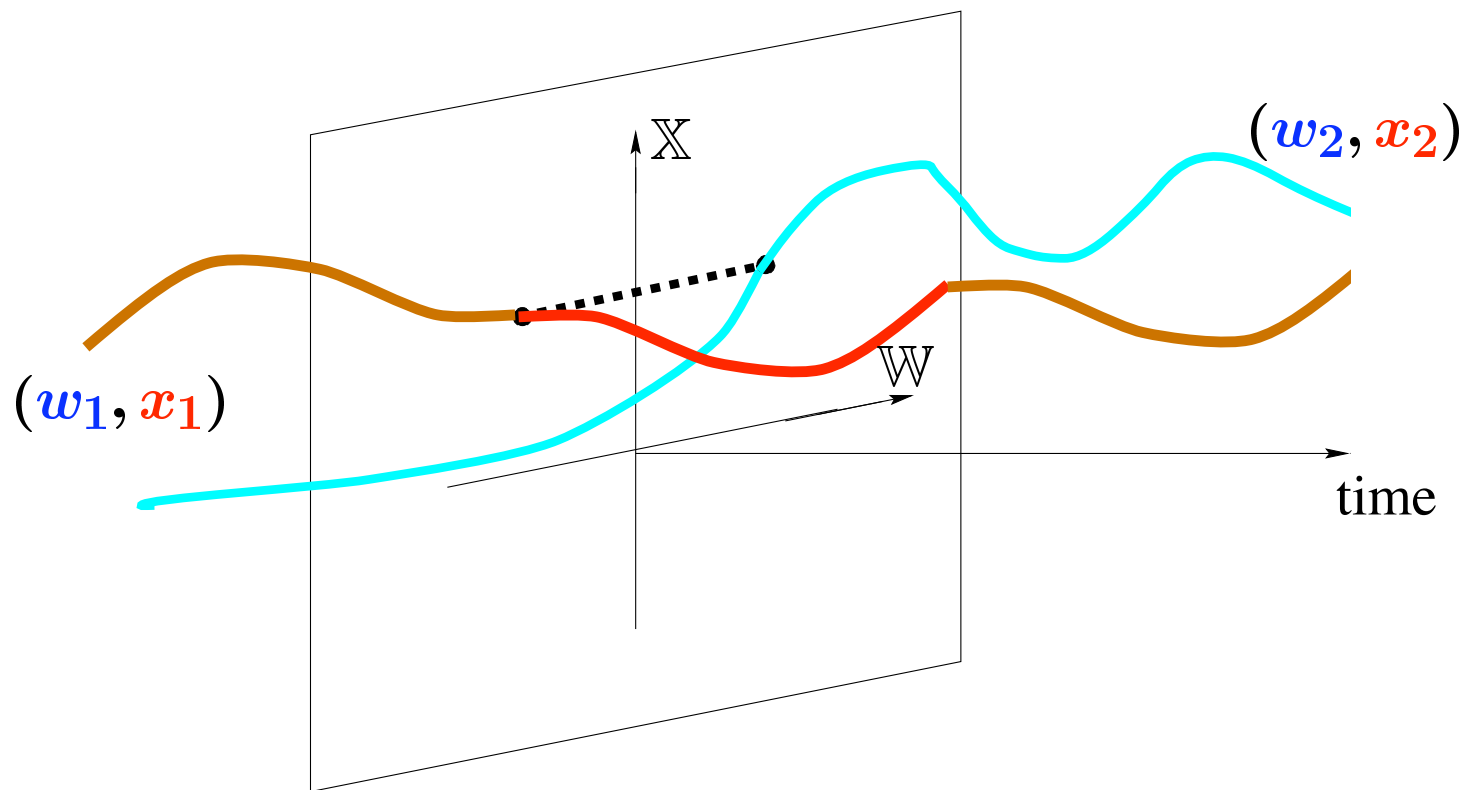
imply

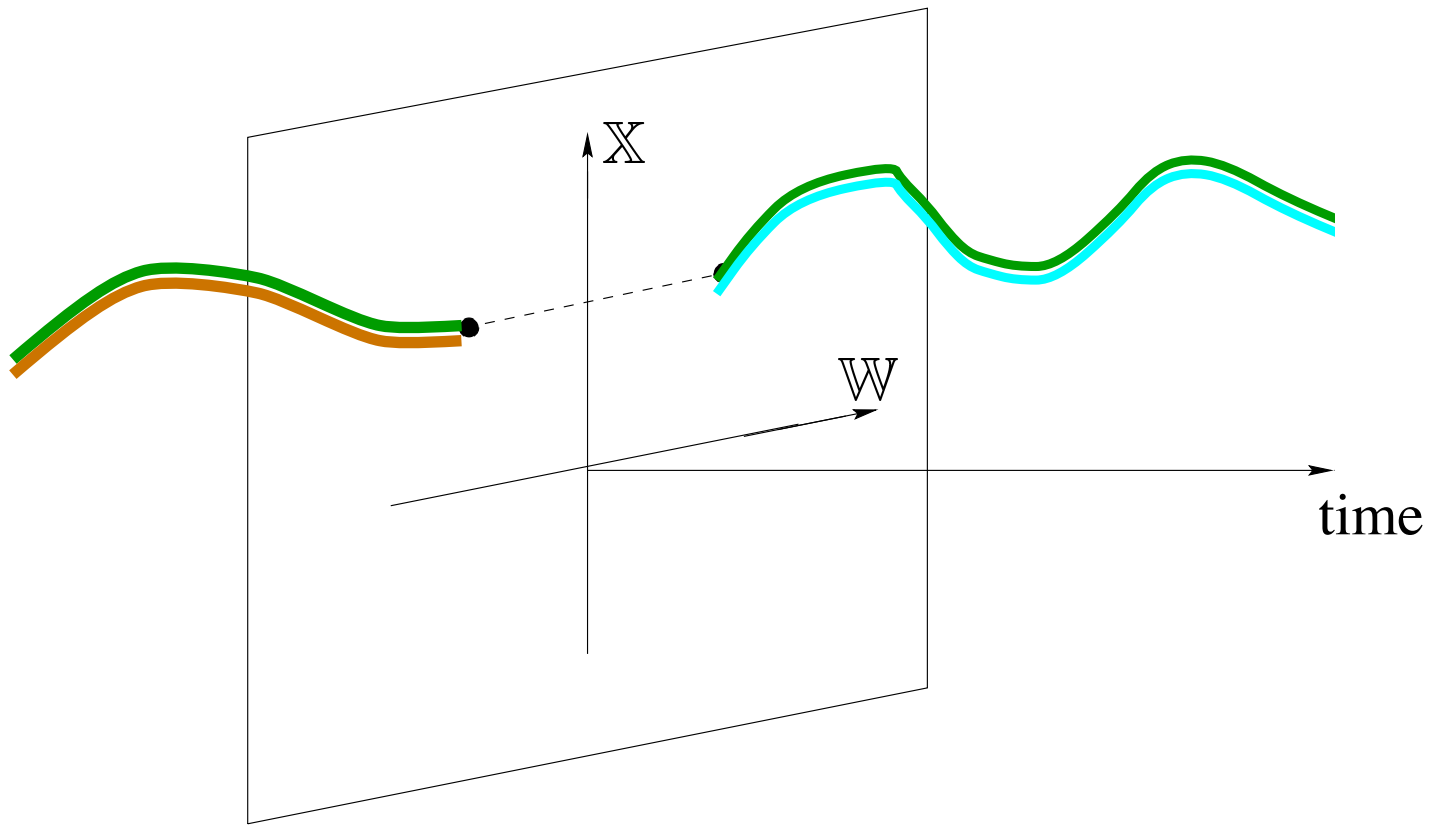
$$(w_1, x_1) \underset{t_0}{\wedge} (w_2, x_2) \in \mathfrak{B}_{\text{full}}.$$

$\underset{t_0}{\wedge}$ denotes **concatenation** at t_0 , defined as

$$f_1 \underset{t_0}{\wedge} f_2(t) := \begin{cases} f_1(t) & \text{for } t < t_0 \\ f_2(t) & \text{for } t \geq t_0 \end{cases}$$

In pictures:





Concatenation

This definition is the implementation of the idea:

The state at time t , $\mathbf{x}(t)$, contains all the information (about $(\mathbf{w}, \mathbf{x})!$) that is relevant for the future behavior.

The state = the **memory (nothing ‘minimal’ implied!).**

**The past and the future are ‘independent’,
conditioned on (given) the present state.**

Examples of state systems:

1. Discrete-time systems.

A latent variable system described by a difference equation that is *first order* in the **latent** variable x , and *zero-th order* in the **manifest** variable w :

$$F(x(t+1), x(t), w(t), t) = 0.$$

2. Continuous-time systems.

A latent variable system described by a differential equation that is *first order* in the **latent** variable \mathbf{x} , and *zero-th order* in the **manifest** variable w :

$$F\left(\frac{d}{dt}\mathbf{x}(t), \mathbf{x}(t), w(t), t\right) = 0.$$

In particular, the ubiquitous

$$\frac{d}{dt}\mathbf{x}(t) = f(\mathbf{x}(t), u(t)), y(t) = h(\mathbf{x}(t), u(t));$$

$$w(t) = (u(t), y(t)).$$

3. Automata.

4. Trellis diagrams.

5. QM:

$$\frac{d}{dt}\psi = i\hbar H(\psi), \quad p = |\psi|^2;$$

ψ = the ‘wave function’;

$p(x, t)$ = the ‘probability’ density of the particle’s position.

For discrete time state systems \rightsquigarrow

Theorem: The ‘complete’ latent variable system

$$\Sigma_X = (\mathbb{Z}, \mathbb{W}, \mathbb{X}, \mathcal{B}_{\text{full}})$$

is a state system if and only if $\mathcal{B}_{\text{full}}$ admits a representation as a difference equation that is *first order* in the **latent** variable x , and *zero-th order* in the **manifest** variable w :

$$F(x(t+1), x(t), w(t), t) = 0.$$

LINEAR DIFFERENTIAL SYSTEMS

We now introduce the systems

$$\Sigma = (\mathbb{R}, \mathbb{R}^w, \mathfrak{B})$$

that are

1. **linear**, meaning
 $((w_1, w_2 \in \mathfrak{B}) \wedge (\alpha, \beta \in \mathbb{R})) \Rightarrow (\alpha w_1 + \beta w_2 \in \mathfrak{B});$
2. **time-invariant**, meaning
 $((w \in \mathfrak{B}) \wedge (t \in \mathbb{R})) \Rightarrow (\sigma^t w \in \mathfrak{B}),$
where σ^t denotes the backwards t -shift;
3. **differential**, meaning
 \mathfrak{B} consists of the solutions of a system of differential equations.

In vector/matrix notation:

$$w = \begin{bmatrix} w_1 \\ w_2 \\ \vdots \\ w_w \end{bmatrix}, \quad R_k = \begin{bmatrix} R_{1,1}^k & R_{1,2}^k & \cdots & R_{1,w}^k \\ R_{2,1}^k & R_{2,2}^k & \cdots & R_{2,w}^k \\ \vdots & \vdots & \cdots & \vdots \\ R_{g,1}^k & R_{g,2}^k & \cdots & R_{g,w}^k \end{bmatrix} \cdot$$

Yields

$$R_0 w + R_1 \frac{d}{dt} w + \cdots + R_n \frac{d^n}{dt^n} w = 0,$$

with $R_0, R_1, \dots, R_n \in \mathbb{R}^{g \times w}$.

Combined with the polynomial matrix

$$R(\xi) = R_0 + R_1\xi + \cdots + R_n\xi^n,$$

we obtain the short notation

$$R\left(\frac{d}{dt}\right)w = 0.$$

Including latent variables \rightsquigarrow

$$R\left(\frac{d}{dt}\right)w = M\left(\frac{d}{dt}\right)\ell$$

with $R, M \in \mathbb{R}^{\bullet \times \bullet}[\xi]$.

Examples:

1. RLC-circuit: Case 1: $CR_C \neq \frac{L}{R_L}$.

Then the relation between V and I is

$$\begin{aligned} \left(\frac{R_C}{R_L} + \left(1 + \frac{R_C}{R_L}\right) CR_C \frac{d}{dt} + CR_C \frac{L}{R_L} \frac{d^2}{dt^2} \right) V \\ = \left(1 + CR_C \frac{d}{dt}\right) \left(1 + \frac{L}{R_L} \frac{d}{dt}\right) R_C I. \end{aligned}$$

We have $w = 2$; $g = 1$; $w = \begin{bmatrix} V \\ I \end{bmatrix}$; $R(\xi) =$

$$\begin{aligned} & \left[\left(\frac{R_C}{R_L} + \left(1 + \frac{R_C}{R_L}\right) CR_C \xi + CR_C \frac{L}{R_L} \xi^2 \mid -1 - \left(CR_C + \frac{L}{R_L}\right) \xi - \left(CR_C \frac{L}{R_L}\right) \xi^2 \right] \\ & = \left[\frac{R_C}{R_L} \mid -1 \right] + \left[1 + \frac{R_C}{R_L} \mid -CR_C - \frac{L}{R_L} \right] \xi + \left[CR_C \frac{L}{R_L} \mid -CR_C \frac{L}{R_L} \right] \xi^2 \end{aligned}$$

2. Linear systems:

- The ubiquitous

$$P\left(\frac{d}{dt}\right)y = Q\left(\frac{d}{dt}\right)u, \quad w = (u, y)$$

with $P, Q \in \mathbb{R}^{\bullet \times \bullet}[\xi]$, $\det(P) \neq 0$ and, perhaps, $P^{-1}Q$ proper.

- The ubiquitous

$$\frac{d}{dt}\mathbf{x} = A\mathbf{x} + B\mathbf{u}; \quad \mathbf{y} = C\mathbf{x} + D\mathbf{u}, \quad w = (u, y).$$

- The descriptor systems

$$\frac{d}{dt}E\mathbf{x} + F\mathbf{x} + G\mathbf{w} = 0.$$

For simplicity of exposition, we assume smooth solutions.

Whence, $R(\frac{d}{dt})w = 0$ defines the system $\Sigma = (\mathbb{R}, \mathbb{R}^w, \mathfrak{B})$ with

$$\mathfrak{B} = \{w \in \mathcal{C}^\infty(\mathbb{R}, \mathbb{R}^w) \mid R(\frac{d}{dt})w = 0\}.$$

NOTATION

\mathcal{L}^\bullet : all such systems (with any - finite - number of variables)

\mathcal{L}^w : with w variables

$\mathcal{B} = \ker(R(\frac{d}{dt}))$

$\mathcal{B} \in \mathcal{L}^w$ (no ambiguity regarding \mathbb{T}, \mathbb{W})

The theory (representations, elimination, controllability, observability, algorithms, control, etc.) of \mathcal{L}^\bullet is very complete.

STATE REPRESENTATIONS

DESCRIPTOR SYSTEMS

Theorem: The latent variable system $(\mathbb{R}, \mathbb{R}^w, \mathbb{R}^n, \mathcal{B}_{\text{full}})$ with $\mathcal{B}_{\text{full}} \in \mathcal{L}^{w+n}$ is a state system if and only if $\mathcal{B}_{\text{full}}$ admits a kernel representation that is *first order* in the **latent** variable \mathbf{x} , and *zero-th order* in the **manifest** variable w .

In other words, iff there exist matrices $E, F, G \in \mathbb{R}^{\bullet \times \bullet}$ such that this kernel representation takes the form of a *descriptor system*:

$$E \frac{d}{dt} \mathbf{x} + F \mathbf{x} + G w = 0.$$

MINIMALITY of STATE REPRESENTATIONS

We can consider two types of **minimality**:

1. Minimality of **the number of equations**
2. Minimality of **the number of state variables**

We discuss mainly the second one.

Definition: The state system $(\mathbb{R}, \mathbb{R}^w, \mathbb{R}^n, \mathcal{B}_{\text{full}})$ with $\mathcal{B}_{\text{full}} \in \mathcal{L}^{w+n}$ is said to be **state-minimal** if, whenever $(\mathbb{R}, \mathbb{R}^w, \mathbb{R}^{n'}, \mathcal{B}'_{\text{full}})$ with $\mathcal{B}'_{\text{full}} \in \mathcal{L}^{w+n'}$ is another state system with the same manifest behavior, there holds

$$n \leq n'.$$

One more definition...

$\mathfrak{B} \in \mathcal{L}^w$ is said to be **trim** if, for all $w_0 \in \mathbb{R}^w$, there exists $w \in \mathfrak{B}$ such that $w(0) = w_0$. The state system $(\mathbb{R}, \mathbb{R}^w, \mathbb{R}^n, \mathfrak{B}_{\text{full}})$ with $\mathfrak{B}_{\text{full}} \in \mathcal{L}^{w+n}$ is said to be **state-trim** if, for all $x_0 \in \mathbb{R}^n$, there exists $(w, x) \in \mathfrak{B}_{\text{full}}$ such that $x(0) = x_0$.

Theorem:

The state system $(\mathbb{R}, \mathbb{R}^w, \mathbb{R}^n, \mathfrak{B}_{\text{full}})$ with $\mathfrak{B}_{\text{full}} \in \mathcal{L}^{w+n}$ is *state-minimal* iff it is **state trim** and the state x is **'observable'** from w .

Observability $:\Leftrightarrow x$ can be deduced from w .

I.e., $\exists X \in \mathbb{R}^{n \times w}[\xi]$ such that $(w, x) \in \mathfrak{B}_{\text{full}} \Leftrightarrow x = X\left(\frac{d}{dt}\right)w$.

State-minimality \Leftrightarrow the combination of **trimness** and **observability**.

Notes:

1. **State isomorphism theorem.** Assume $(\mathbb{R}, \mathbb{R}^w, \mathbb{R}^n, \mathfrak{B}_{\text{full}})$ and $(\mathbb{R}, \mathbb{R}^w, \mathbb{R}^n, \mathfrak{B}'_{\text{full}})$, $\mathfrak{B}_{\text{full}}, \mathfrak{B}'_{\text{full}} \in \mathcal{L}^{w+n}$ both state-minimal, same manifest behavior
 \Rightarrow there exists a nonsingular $S \in \mathbb{R}^{n \times n}$ such that
 $((w, x) \in \mathfrak{B}_{\text{full}} \text{ and } (w, x') \in \mathfrak{B}'_{\text{full}}) \Leftrightarrow (x' = Sx).$
2. The manifest behavior is **controllable** iff a state-minimal state representation is **state-controllable** (defined in the obvious way).
3. \exists algorithms on E, F, G in a descriptor representation to verify its state-minimality, its equation minimality, both combined.

4. $E \frac{d}{dt}x + Fx + Gw = 0$ and $E' \frac{d}{dt}x' + F'x' + G'w = 0$ are two minimal (state- and equation-minimal) representations of the same manifest behavior iff there exist nonsingular matrices $T, S \in \mathbb{R}^{\bullet \times \bullet}$ such that

$$E' = TES, F' = TES, G' = TG.$$

All 'classical' results remain valid, except, (fortunately!)
the celebrated (non-)equivalence:

state-minimality \Leftrightarrow (state-observability + state-controllability).

Non-controllable systems are **very 'real'** and they allow state-minimal (non-controllable) state representation.

STATE CONSTRUCTION

**Given a representation of the manifest behavior,
find a (state-minimal) state representation for it.**

**Most logical : latent variable representation \rightsquigarrow state representation.
However, it is most convenient to discuss kernel representations first.**

Let $X(\xi) \in \mathbb{R}^{\bullet \times w}[\xi]$. The map $X\left(\frac{d}{dt}\right)$ is called a **state map** for $\mathfrak{B} \in \mathcal{L}^w$ if the full behavior

$$\mathfrak{B}_{\text{full}} = \left\{ (w, x) \mid w \in \mathfrak{B} \text{ and } x = X\left(\frac{d}{dt}\right)w \right\}$$

satisfies the axiom of state. **Minimal state map:** obvious.

In a state-minimal representation, x is always determined by a state map (because of observability), whence state maps exist.

CONSTRUCTION of STATE MAPS

Define the *'shift-and-cut'* operator σ on $\mathbb{R}[\xi]$ as follows:

$$\begin{aligned}\sigma : p_0 + p_1\xi + \cdots + p_{n-1}\xi^{n-1} + p_n\xi^n \\ \mapsto p_1 + p_2\xi + \cdots + p_{n-1}\xi^{n-2} + p_n\xi^{n-1}\end{aligned}$$

Extend-able in the obvious term-by-term way to $\mathbb{R}^{\bullet \times \bullet}[\xi]$.

Repeated use of the cut-and-shift on $P \in \mathbb{R}^{\bullet \times \bullet}$ yields the *'stack'* operator Σ_P , defined by

$$\Sigma_P := \begin{bmatrix} \sigma(P) \\ \sigma^2(P) \\ \vdots \\ \sigma^{\text{degree}(P)}(P) \end{bmatrix}$$

FROM KERNEL REPRESENTATION to STATE MAP

There is a construction (elegant in its simplicity) of a state map in terms of the cut-and-shift and stack operators!

Theorem: Let $R(\frac{d}{dt})w = 0$ be a kernel representation of $\mathfrak{B} \in \mathcal{L}^w$. Then $\Sigma_R(\frac{d}{dt})$ is a state map for \mathfrak{B} .

The resulting state representation

$$R(\frac{d}{dt})w = 0 ; \quad x = \Sigma_R(\frac{d}{dt})w$$

need not be minimal. It is trivially state-observable, but it may not be state-trim. Using **Gröbner basis techniques** it can be trimmed, leading to a minimal state representation.

SINGLE INPUT - SINGLE OUTPUT SYSTEMS

Apply this to

$$p\left(\frac{d}{dt}\right)y = q\left(\frac{d}{dt}\right)u$$

with

$$p(\xi) = p_0 + p_1\xi + \cdots + p_{n-1}\xi^{n-1} + p_n\xi^n, \quad p_n \neq 0$$

$$q(\xi) = q_0 + q_1\xi + \cdots + q_{n-1}\xi^{n-1} + q_n\xi^n$$

The cut-and-shift and stack operators yield the polynomial matrix

$$X(\xi) = \begin{bmatrix} p_1 + \cdots + p_{n-1}\xi^{n-2} + p_n\xi^{n-1} & -q_1 - \cdots - q_{n-1}\xi^{n-2} - q_n\xi^{n-1} \\ p_2 + \cdots + p_{n-1}\xi^{n-3} + p_n\xi^{n-2} & -q_2 - \cdots - q_{n-1}\xi^{n-3} - q_n\xi^{n-2} \\ \vdots & \vdots \\ p_{n-1} + p_n\xi & -q_{n-1} - q_n\xi \\ p_n & -q_n \end{bmatrix}$$

It follows that $x = X\left(\frac{d}{dt}\right)$ is a state map, in fact, a **state minimal** one, even if the system is not controllable, i.e., if p and q have a common factor.

Minimal state map \cong basis for span of the rows of X modulo R .

Number the rows of X in reverse order. A small calculation shows that this choice of the state variables leads to the so-called

observer canonical form, the i/s/o representation

$$A = \begin{bmatrix} -p_0/p_n & 1 & 0 & \cdots & 0 & 0 \\ -p_1/p_n & 0 & 1 & \cdots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ -p_{n-1}/p_n & 0 & 0 & \cdots & 0 & 1 \end{bmatrix}, \quad B = \begin{bmatrix} q_0 - p_0 q_n/p_n \\ q_1 - p_1 q_n/p_n \\ \vdots \\ q_{n-1} - p_{n-1} q_n/p_n \end{bmatrix},$$

$$C = [1/p_n \ 0 \ 0 \ \cdots \ 0 \ 0], \quad D = [q_n/p_n].$$

Another immediate choice is to pick linear combinations of the rows of X so that the resulting state map matrix takes the form

$$X'(\xi) = \begin{bmatrix} 1 & \star \\ \xi & \star \\ \vdots & \vdots \\ \xi^{n-2} & \star \\ \xi^{n-1} & \star \end{bmatrix}$$

The \star 's are obtained from the polynomial $b(\xi) \in \mathbb{R}[\xi]$ defined by the equation

$$p(\xi)b(\xi^{-1}) = q(\xi) \quad (\text{modulo } \xi^{-1}\mathbb{R}[\xi^{-1}]).$$

Then

$$X'(\xi) = \begin{bmatrix} 1 & b_0 \\ \xi & b_1 + b_0\xi \\ \vdots & \vdots \\ \xi^{n-2} & b_{n-2} + b_{n-3}\xi + \cdots + b_0\xi^{n-2} \\ \xi^{n-1} & b_{n-1} + b_{n-2}\xi + \cdots + b_0\xi^{n-1} \end{bmatrix}.$$

This leads to the **observable canonical form**, the i/s/o representation

$$A = \begin{bmatrix} 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 1 \\ -\frac{p_0}{p_n} & -\frac{p_1}{p_n} & -\frac{p_2}{p_n} & \cdots & -\frac{p_{n-1}}{p_n} \end{bmatrix}, \quad B = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_{n-1} \\ b_n \end{bmatrix},$$

$$C = [1 \ 0 \ \cdots \ 0 \ 0], \quad D = [b_0].$$

Image representations + state construction \rightsquigarrow

controller canonical form and **controllable canonical form**.

Notes:

- **Basic idea of algorithms:**
from latent variable representation directly to state model.
This complements the existent algorithms
transfer function \rightarrow i / s / o representation;
impulse response \rightarrow i / s / o representation.
- **Our state construction is easily extended to state / input construction.**

- **Examples of useful special (minimal) state representations:**

i/s/o representation:

$$\frac{d}{dt}x = Ax + Bu, \quad y = Cx + Du, \quad w = (u, y),$$

output nulling representation:

$$\frac{d}{dt}x = Ax + Bw, \quad 0 = Cx + Dw,$$

driving variable representation:

$$\frac{d}{dt}x = Ax + Bv, \quad w = Cx + Dv.$$

- **Immediately deduced from descriptor representation:**

$$E \frac{d}{dt}x + Fx + Gw = 0.$$

BALANCED REALIZATIONS

FROM KERNEL REPR. TO BALANCED REAL.

For simplicity, we discuss only single input - single output systems

$$p\left(\frac{d}{dt}\right)y = q\left(\frac{d}{dt}\right)u$$

assumed controllable, i.e., with p, q co-prime,
 $n = \text{degree}(p) \geq \text{degree}(q)$,
and stable, i.e., p Hurwitz.

Problem: Pass from p, q to a balanced state representation using polynomial algebra.

**Computation of the controllability and observability Gramians
(two-variable polynomials):**

$$C(\zeta, \eta) = \frac{p(\zeta)p(\eta) - p(-\zeta)p(-\eta)}{\zeta + \eta},$$

$$O(\zeta, \eta) = \frac{q(\zeta)q(\eta) - x(\zeta)p(\eta) - p(\zeta)x(\eta)}{\zeta + \eta},$$

where $x \in \mathbb{R}[\xi]$ is given by the Bezout equation

$$x(-\xi)p(\xi) + p(-\xi)x(\xi) = q(-\xi)q(\xi).$$

Factor C, O as

$$C(\zeta, \eta) = \sum_{k=1}^n \sigma_k^{-1} x_k(\zeta) x_k(\eta),$$

$$O(\zeta, \eta) = \sum_{k=1}^n \sigma_k x_k(\zeta) x_k(\eta),$$

with $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_n > 0$. Solve for $k = 1, \dots, n$, the equations

$$\xi x_k(\xi) = \sum_{k'=1}^n a_{k,k'} x_{k'}(\xi) + b_k p(\xi),$$

$$q(\xi) = \sum_{k'=1}^n c_{k'} x_{k'}(\xi) + d p(\xi).$$

Then the matrices with entries $a_{k,k'}, b_k, c_k, d$ define a **balanced realization**. The balanced reduction is obvious from here.

Open Problem:

Algorithm

$$(p, q) \rightsquigarrow (p_{\text{balanced}}, q_{\text{balanced}}).$$

AAK version of this: cfr. Fuhrmann's book.

Goal: algorithms passing from image, latent variable representation to balanced realization.

STATE for PDE's

BEHAVIORS of n-D SYSTEMS

A system := $\Sigma = (\mathbb{T}, \mathbb{W}, \mathcal{B})$

\mathbb{T} = the set of independent variables
time, space, time and space

\mathbb{W} = the set of dependent variables (= where the variables take on their values), signal space, space of field variables, ...

$\mathcal{B} \subseteq \mathbb{W}^{\mathbb{T}}$: the behavior = the admissible trajectories

$$\Sigma = (\mathbb{T}, \mathbb{W}, \mathfrak{B})$$

for a trajectory $w : \mathbb{T} \rightarrow \mathbb{W}$, we thus have:

$w \in \mathfrak{B}$: the model **allows** the trajectory w ,

$w \notin \mathfrak{B}$: the model **forbids** the trajectory w .

In this section, $\mathbb{T} = \mathbb{R}^n$, ('n-D systems')

$$\mathbb{W} = \mathbb{R}^w,$$

$$w : \mathbb{R}^n \rightarrow \mathbb{R}^w, (w_1(x_1, \dots, x_n), \dots, w_w(x_1, \dots, x_n)),$$

often, $n = 4$, independent variables (t, x, y, z) ,

\mathfrak{B} = solutions of a system of constant coefficient
linear PDE's.

'Linear distributed differential systems'.

n-D LINEAR DIFFERENTIAL SYSTEMS

$$T = \mathbb{R}^n,$$

$$W = \mathbb{R}^w,$$

\mathfrak{B} = the solutions of a linear constant coefficient system of PDE's.

Let $R \in \mathbb{R}^{\bullet \times w}[\xi_1, \dots, \xi_n]$, and consider

$$R\left(\frac{\partial}{\partial x_1}, \dots, \frac{\partial}{\partial x_n}\right)w = 0 \quad (*)$$

Define

$$\mathfrak{B} := \{w \in \mathcal{C}^\infty(\mathbb{R}^n, \mathbb{R}^w) \mid (*) \text{ holds} \}$$

$\mathcal{C}^\infty(\mathbb{R}^n, \mathbb{R}^w)$ **mainly** for convenience.

Example: Maxwell's equations

$$\begin{aligned}\nabla \cdot \vec{E} &= \frac{1}{\epsilon_0} \rho, \\ \nabla \times \vec{E} &= -\frac{\partial}{\partial t} \vec{B}, \\ \nabla \cdot \vec{B} &= 0, \\ c^2 \nabla \times \vec{B} &= \frac{1}{\epsilon_0} \vec{j} + \frac{\partial}{\partial t} \vec{E}.\end{aligned}$$

$T = \mathbb{R} \times \mathbb{R}^3$ (time and space),

$w = (\vec{E}, \vec{B}, \vec{j}, \rho)$

(electric field, magnetic field, current density, charge density),

$W = \mathbb{R}^3 \times \mathbb{R}^3 \times \mathbb{R}^3 \times \mathbb{R}$,

$\mathfrak{B} =$ set of solutions to these PDE's.

Note: 10 variables, 8 equations! $\Rightarrow \exists$ free variables.

Maxwell's equation for the electrical variables (with \vec{B} 'eliminated'):

$$\begin{aligned}\nabla \cdot \vec{E} &= \frac{1}{\epsilon_0} \rho, \\ \epsilon_0 \frac{\partial}{\partial t} \nabla \cdot \vec{E} + \nabla \cdot \vec{j} &= 0, \\ \epsilon_0 \frac{\partial^2}{\partial t^2} \vec{E} + \epsilon_0 c^2 \nabla \times \nabla \times \vec{E} + \frac{\partial}{\partial t} \vec{j} &= 0.\end{aligned}$$

MARKOVIAN n-D SYSTEMS

What is the notion of state for such systems?

What does 'Markov' mean?

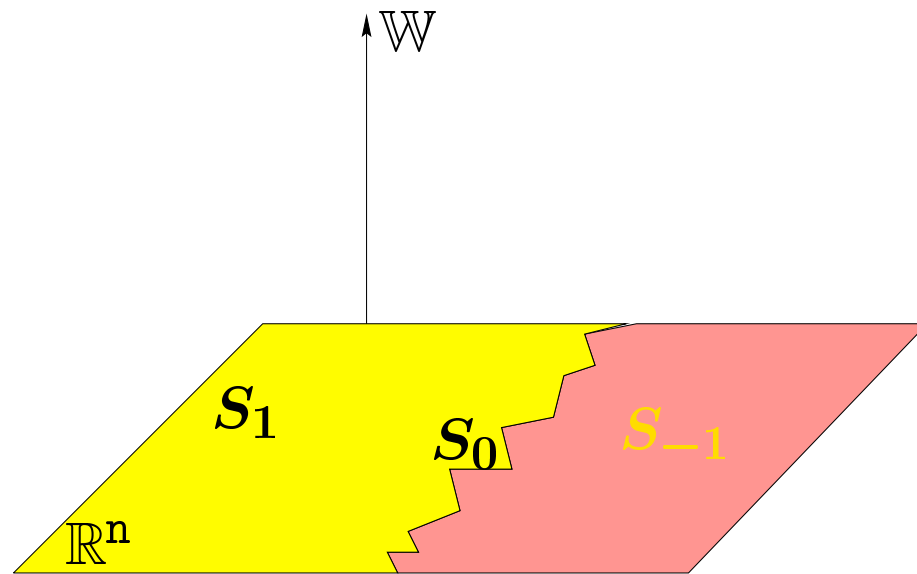
$\mathfrak{B} \in \mathfrak{L}_n^w$ is said to be **Markovian** if for all **nice** partitions of $\mathbb{R}^n = S_{-1} \cup S_0 \cup S_1$ there holds:

$$w_1, w_2 \in \mathfrak{B} \quad \text{and} \quad w_1|_{S_0} = w_2|_{S_0}$$

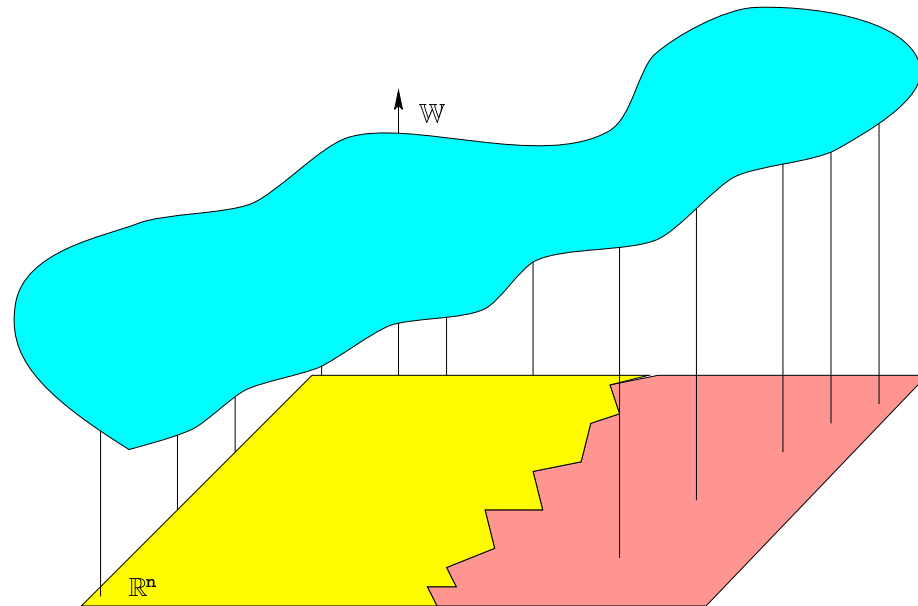
\implies

$$w_1 \underset{S_0}{\wedge} w_2 \in \mathfrak{B}.$$

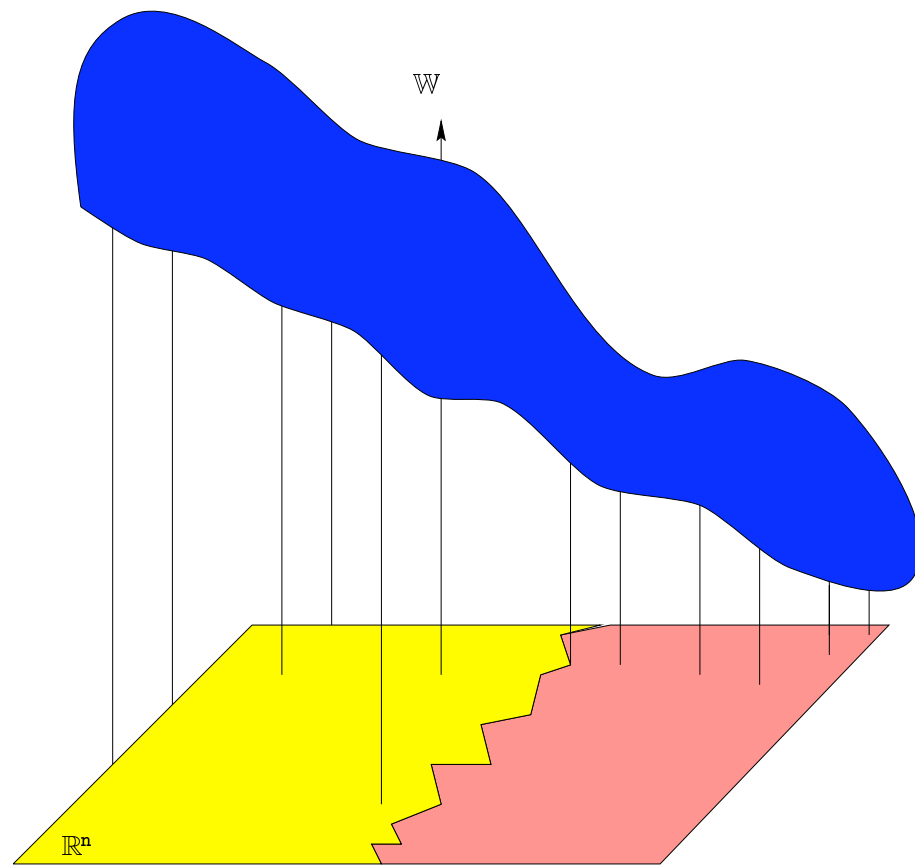
($\underset{S_0}{\wedge}$:= 'concatenation' at S_0).



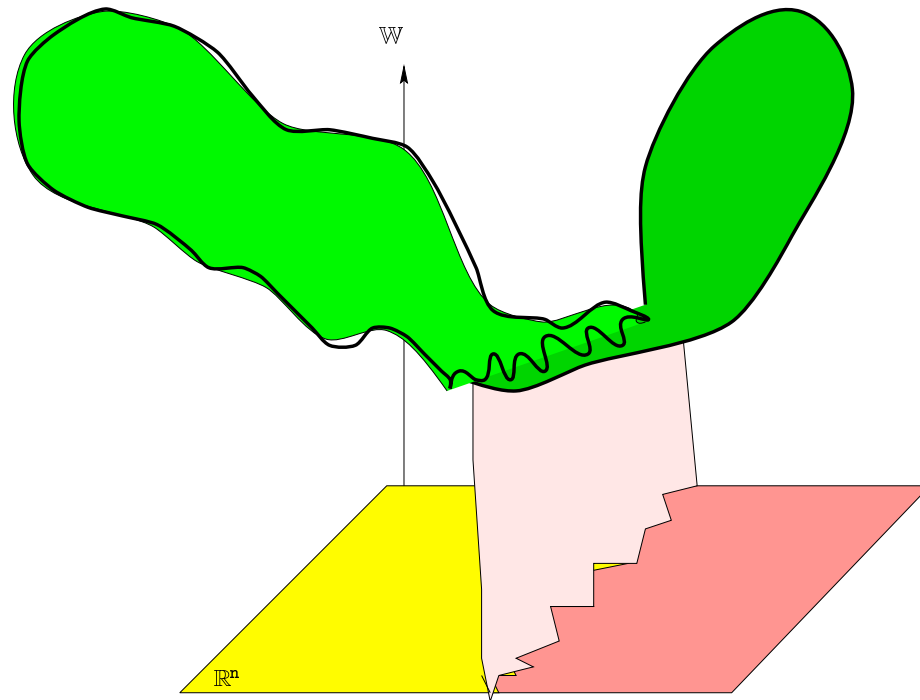
A nice partition



The first solution



The second solution



The concatenated solution

CONJECTURE

$\mathfrak{B} \in \mathcal{L}_n^w$ is Markovian if and only if

$$\mathfrak{B} = \ker\left(R\left(\frac{\partial}{\partial x_1}, \dots, \frac{\partial}{\partial x_n}\right)\right),$$

with R first order, i.e.,

$$R(\xi_1, \dots, \xi_n) = R_0 + R_{1,1}\xi_1 + R_{1,2}\xi_2 + \dots + R_{1,n}\xi_n.$$

“If”-part is clear; “only if”-part is the problem.

Example: Maxwell’s equations, \vec{B} induces state representation of electrical behavior. **Not observable**, thou.

CONCLUDING REMARKS

- A system \cong a **manifest behavior**
- First principles models \rightsquigarrow systems with **latent variables**
- **State systems**: latent variable systems in which the state 'splits' the past and the future
- State construction: for linear differential systems via **cut-and-shift map** and **Gröbner basis algorithms**
- **Balanced reduction** via polynomial algebra
- Conjecture for PDE's: **Markovian** \Leftrightarrow **first order**

Thank you!